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### Investigating the robustness of adaptive dynamic loop scheduling on heterogeneous computing systems

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## Outline

# → Motivation & approach

Contribution Background & related work Robustness metrics design Usefulness & cost analysis Conclusions & future work







## Motivation



- New classes of applications: (very) large & complex, computationally intensive, irregular behavior, contain large loops
- + High performance computing systems: unstable environments, complex to manage, computing resources vary in type, quantity and availability → multiple sources and types of uncertainty







### Motivation



+ Performance is more than just execution time:

+ Scalability

+Execution time, efficiency and other metrics

Numerical efficiency

+Accuracy, stability







### Motivation

- Heterogeneity in computing resources evolves from variations in:
  - Number → failures (fault tolerance issues)
  - Load → availability (load balancing issues)







## Challenges



User goals: **unchanged:** optimal results with minimum cost/effort

Pradeoff between numerical efficiency and scalability

Tradeoff between scheduling overhead and load imbalance



How to execute large scientific applications on today's heterogeneous systems in a **flexible** and **resilient**, taken together as *robustness*, manner?



















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**Contribution: flexibility** and **resilience metrics** to quantify the **robustness** (reliability) of 3 Adaptive **DLS** (ADLS) methods against *load variations* and *resource failures*, respectively

#### We propose:

Use of the FePIA procedure to design metrics to measure robustness of batched AWF, chunked AWF, and AF to schedule large loop iterations in a **flexible** and **resilient** manner on large-scale heterogeneous computing systems







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### Background & related work

Bobustness encompasses *flexibility and resilience* 

Previous study shows it is feasible to develop a common metric for all non-adaptive and one adaptive DLS techniques

#### What is missing?

A more general approach needed to address robustness for more than one method (DLS or ADLS) and for more than one application







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### Robustness metrics - FePIA Design steps



- **S.1** Identify the set of performance features of the DLS method ( $\Phi$ )
- S.2 Identify the set of parameters perturbing the performance of the method ( $\Pi$ )
- **5.3** Identify and clarify the impact of the perturbation parameters on the performance features of the DLS method ( $\Phi(\Pi)$ )
- **5.4** Identify the analysis to determine the robustness of the DLS method  $(r(\varphi, \pi) robustness radius, \rho(\Phi, \Pi) robustness metric)$

We use these steps to design the following metrics:  $\rho(\Phi_1, \Pi_1)$ : Robustness against *load variations* (flexibility)  $\rho(\Phi_2, \Pi_2)$ : Robustness against *resource failures* (resilience)







### **Robustness metrics - Notation**

Notation	Definition
N	total number of tasks
$N^{resch}$	# of tasks that need to be <i>rescheduled</i>
$a_i$	<i>i</i> -th task, $1 \leq i \leq N$
P	total number of processors
$m_j$	<i>j</i> -th processor, $1 \leq j \leq P$
$T_{j}$	<i>execution</i> time of task $a_i$ on $m_j$
$T_{ij}^{W2F}$	<i>communication</i> time between $m_j$ and foreman for $a_i$
$T_{ii}^{W2W}$	<i>communication</i> time between $m_i$ and any other worker for $a_i$
εŤi	<i>finishing</i> time of all tasks computed by $m_i$
$T_{PAR}$	total parallel execution time for the N tasks
$T_{c rot}^{fastest}$	execution time of the fastest sequential version of the application
3 EQ	$_T fastest$
$S_p$	Speed up of the parallel system defined as $\frac{-SEQ}{TPAP}$
$\lambda = [\lambda_1 \dots \lambda_P]^T$	vector of processors <i>load</i> (= system load)
$\mathbf{F} = [\mathbf{f}_1 \dots \mathbf{F}_{\mathbf{P}}]^T$	resources status vector (active/failed)
$\Phi = \{\phi_1, \dots\}$	set of <i>performance</i> features
$\Pi = \{\pi_1, \dots\}$	set of <i>perturbation</i> parameters
$T_1, T_2, T_3$	tolerance factors for performance features
$r_{DLS}(,)$	robustness radius
$\rho_{DLS}(,)$	robustness metric
W	normalized weight of processors in WF, AWF, AWF-B, AWF-C

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#### Flexibility: robustness w.r.t load variations

Assumptions regarding variations that can occur in the load at run-time:

A.1 variations of individual worker loads

are mutually independent

A.2 individual worker load variations may or may

not occur simultaneously

A.3 ADLS have load variation detection &

monitoring mechanisms



These are realistic assumptions that simplify the cause-andeffect analysis for this metric.







#### Flexibility: robustness w.r.t load variations

- **S.1** Performance features set:  $\Phi_1 = {\phi_1} = {\mathcal{E}T_j}$
- **S.2** Perturbation parameters set:  $\Pi_1 = {\pi_1} = {\Lambda_j}$
- S.3 Impact of  $\Pi_1$  on  $\Phi_1$ : for all {tasks i | i executed on  $m_j$  in presence of  $\Lambda_j$ }

$$\mathcal{E}\mathsf{T}_{j}(\Lambda_{j}) = \Sigma_{i,j}^{N,P}(\mathsf{T}_{ij}(\Lambda_{j}) + \mathsf{T}_{ij}^{w2f}(\Lambda_{j}) + \mathsf{T}_{ij}^{w2w}(\Lambda_{j}))$$

**5.4** Analyze the robustness radius and define the robustness metric:

$$\{ \Lambda_j \in \langle \Lambda_j', \Lambda_j^{''} \rangle \mid (\mathcal{E}\mathsf{T}_j \ (\Lambda_j) = \mathsf{T}_1 \cdot \mathcal{E}\mathsf{T}_j^{\mathsf{orig}}) \land (1 \le j \le \mathsf{P}) \}, \Lambda_j^{\mathsf{orig}} - \text{initial load on } \mathsf{m}_j$$

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#### Flexibility: robustness w.r.t load variations



 $\tau_1$  – acceptable tolerance on  ${\sf ET}_j$  or  ${\sf T}_{\sf PAR}$  impact caused by variations in  $\Lambda_j$  or  $\Lambda$ 

Robustness radii 
$$r_{ADLS}(\varphi_i, \pi_i) = r_{ADLS}(\varepsilon T_j, \Lambda_j) = \max ||\Lambda_j - \Lambda_j^{orig}||_1$$
, s.t.  $\varepsilon T_j(\Lambda_j) = \tau_1 \cdot \varepsilon \tau_j^{orig}$ 

Individual robustness value  $\rho_{ADLS}(\Phi_1, \Pi_1) = \max(r(\mathcal{E}T_j, \Lambda_j))$  $\forall \phi_i \in \Phi_1 \text{ and } \pi_i \in \Pi_1$ 

Robustness metric  $\rho_{metric}(T_{PAR}, \Lambda) = min(\rho_{ADLS}(\epsilon T_j, \Lambda_j))$ 







# Possible scenarios to determine the flexibility of ADLS techniques



Scenario (a) Choose the ADLS method that has the *lowest impact on ADLS* performance AND can handle the *largest variation in*  $\Lambda$ 

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Scenario (b) Choose the ADLS method that has the *lowest impact on ADLS* performance for a *fixed variation in*  $\Lambda$ 

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#### Resilience: robustness w.r.t resource failures

Assumptions regarding failures that can occur in the system and their handling mechanisms:

A.4 only resources associated with worker

processors fail

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- A.5 resource failures occur simultaneously
- A.6 resource failures are mutually independent
- A.7 resource failures are permanent
- A.8 ADLS have fault discovery & fault-recovery mechanisms
- <sup>1</sup> These are realistic assumptions that simplify the cause-and-effect analysis for this metric.





#### Resilience: robustness w.r.t resource failures



- **S.1** Performance features set:  $\Phi_2 = \{\phi'_1, \phi'_2\} = \{N^{\text{resch}}, T_{\text{PAR}}\}$
- **S.2** Perturbation parameters set  $\Pi_2 = {\pi'_1} = {F}$
- **5.3** Impact of  $\Pi_2$  on  $\Phi_2$ :
- $\varphi'_{1}=f_{11}(\pi'_{1}): N^{resch}=f_{11}(\mathbf{F}), \text{ where } N^{resch}(\mathbf{F}) = N_{p}^{resch}(\mathbf{F}) + N_{lb}^{resch}(\mathbf{F}) \land N^{resch}(\mathbf{F}) = f'(ADLS \text{ method of choice})$
- $\varphi'_{2}=f_{21}(\pi'_{1}): T_{PAR}=f_{21}(F)$ , where  $T_{PAR}(F) = f''(ADLS \text{ method & recovery mech.})$
- **5.4** Analyze the robustness radius/radii and define the robustness metric:

 $\{ \mathbf{F} \mid (N^{\text{resch}}(\mathbf{F}) \leq \mathbf{T_2} \cdot \mathbf{N}) \land (\exists \mathbf{F}' \text{ s.t. } N^{\text{resch}}(\mathbf{F}') > \mathbf{T_2} \cdot \mathbf{N}) \}$ 

 $\{ \mathbf{F} \mid (\mathsf{T}_{\mathsf{PAR}}(\mathbf{F}) \leq \mathbf{T}_{\mathbf{3}} \cdot \mathsf{T}_{\mathsf{PAR}}(\mathbf{F}^{\mathsf{orig}})) \land (\exists \mathbf{F}' \mathsf{s.t.} \mathsf{T}_{\mathsf{PAR}}(\mathbf{F}') > \mathbf{T}_{\mathbf{3}} \cdot \mathsf{T}_{\mathsf{PAR}}(\mathbf{F}') \}$ 







### Resilience: robustness w.r.t resource failures (cont.)

Robustness radii (tolerance intervals):

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 $\begin{aligned} r_{ADLS}(\varphi'_{1},\pi'_{1})=r_{ADLS}(N^{resch},F) &= max||F-F^{orig}||_{1}, \text{ such that} \\ (N^{resch}(F)\leq \tau_{2}\cdot N) \land (\exists F' s.t. N^{resch}(F')>\tau_{2}\cdot N) \end{aligned}$ 

 $\begin{aligned} r_{ADLS}\left(\phi_{2}^{\prime},\pi_{1}^{\prime}\right)=r_{ADLS}\left(\mathsf{T}_{PAR},\mathsf{F}\right)=\max\left|\left|\mathsf{F}-\mathsf{F}^{\operatorname{orig}}\right|\right|_{1}, \text{ such that}\\ \left(\mathsf{T}_{PAR}(\mathsf{F})\leq_{\mathbf{T}_{3}}\cdot\mathsf{T}_{PAR}\left(\mathsf{F}^{\operatorname{orig}}\right)\right)\wedge\left(\exists \mathsf{F}^{\prime}\mathsf{s.t.} \mathsf{T}_{PAR}(\mathsf{F}^{\prime})>_{\mathbf{T}_{3}}\cdot\mathsf{T}_{PAR}\left(\mathsf{F}^{\prime\operatorname{orig}}\right)\right)\end{aligned}$ 

 $\tau_2$  - acceptable tolerance on N<sup>resch</sup> impact caused by # of failures in **F**.

 $T_3$  - acceptable tolerance on  $T_{PAR}$  impact caused by # of failures in **F**.

Individual robustness value  $\rho_{ADLS}(\Phi_2, \Pi_2) = \max(r_{ADLS}(\phi'_1, \pi'_1), r_{ADLS}(\phi'_2, \pi'_1))$ Robustness metric

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$$\rho_{\text{metric}}(\mathsf{T}_{\text{PAR}},\mathsf{F}) = \min(\rho_{\text{ADLS}}(\mathsf{N}^{\text{resch}},\mathsf{F}),\rho_{\text{ADLS}}(\mathsf{T}_{\text{PAR}},\mathsf{F}))$$
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#### Possible scenarios to determine the resilience of ADLS techniques



Scenario (a) Choose the ADLS method that has the *lowest impact on ADLS* performance AND can handle the *largest variation in* **F** 

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Scenario (b) Choose the ADLS method that has the *lowest impact on ADLS* performance for a *fixed variation in* **F** 

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### Usefulness



- The usefulness of proposed metrics is twofold:
  - the metrics can be formulated offline with application/ system/ algorithm-specific initial values and integrated into the master to guide and adapt autonomously the scheduling decisions
  - the metrics are usable in conjunction with other desired performance metrics (e.g., makespan) for differentiating among DLS that have similar performance w.r.t makespan
- Tolerance factors (i.e.,  $T_1$ ,  $T_2$ ,  $T_3$ ) must accurately reflect real life conditions and realistic scenarios





# Usefulness



Tolerance	Depends	Best	Worst	Average
factor	on	case	case	case
$\tau_1$	Application type	1.0	1.5	1.25
$ au_2$	ADLS method	0% of $N$	50% of $N$	25% of N
	of choice			
$ au_3$	ADLS method	$S_p^{ideal}$	1	$\frac{S_p}{2}$
	of choice,	-	(no $S_p$ )	-
	# of failures,			
	fault detection &			
	recovery mechanism			

A careful choice of the tolerance factors and incorporation into the adaptive DLS methods, renders the proposed metrics useful towards producing *efficient*, *qualitative* and *reliable* schedules for execution of large and complex scientific applications







### Usefulness



- When considering a set of ADLS and DLS techniques one should select as *robustness metric* (p<sub>metric</sub>) the one that gives the smallest robustness radius among all techniques.
- In order to determine the most robust algorithm one should select the DLS or the ADLS technique with the highest robustness value ( $\rho_{ADLS}^{max} = k \cdot \rho_{metric}$ ) as derived from the selected robustness metric.







# Cost analysis



The computational complexity of each metric is driven by the computational complexity of calculating robustness radii for each metric

General optimization problem for robustness: Maximize  $\sum_{k} ||\pi_{k} - \pi_{k}^{orig}||_{p}$ , k>0 subject to  $\beta^{min} \leq \{\Phi = f(\pi_{k})\} \leq \beta^{max}$ where:  $||\pi_{k}||_{p}$  is the  $\mathcal{L}_{p}$ -norm of the perturbation parameter  $\pi_{\kappa}$   $\langle \beta^{min}, \beta^{max} \rangle$  is the tolerance interval,  $\Phi = f(\pi_{k})$  is equivalent to  $\Phi(\pi_{k}) = \tau \cdot \Phi(\pi_{k}^{orig})$  $\tau$  is the tolerance factor of the performance feature

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## Cost analysis

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Optimization problem for **flexibility**:

maximize  $||\Lambda - \Lambda^{orig}||_1$  subject to

 $\mathsf{T}_{\mathsf{PAR}}(\Lambda) \leq \mathsf{T}_1 \cdot \mathsf{T}_{\mathsf{PAR}}(\Lambda^{\mathsf{orig}}) \land \mathsf{T}_{\mathsf{PAR}}(\Lambda) \leq \mathsf{T}^{\mathsf{fastest}}_{\mathsf{SEQ}}$ 

Optimization problem for **resilience**:

maximize ||F-F<sup>orig</sup>||<sub>1</sub> subject to

 $N^{\text{resch}}(F) \leq \tau_2 \cdot N \land \ T_{\text{PAR}}(F) \leq \tau_3 \cdot T_{\text{PAR}}(F^{\text{orig}}) \land \ T_{\text{PAR}}(F) \leq T^{\text{fastest}}_{\text{SEQ}}$ 

**Fact:** All norms  $|| \wedge ||_p$  are convex functions

**Dilemma:**  $\Phi(\Pi)$ : convex or concave function?

If  $\Phi(\Lambda)$  and  $\Phi(F)$  are linear or convex functions both metrics become convex optimization problems with inexpensive optimal solutions

If  $\Phi(\Lambda)$  and  $\Phi(F)$  are concave then near-optimal solutions cannot be found with optimal costs





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### **Conclusions & future work**



C and AF against load variations

The proposed metrics bring out the <u>most adaptive</u> and efficient DLS algorithms to state-of-the-art performance

© Careful choice of tolerance values results in <u>qualitative</u>, <u>efficient</u> and <u>reliable</u> schedules in today's large scale, high-performance, heterogeneous systems







### Conclusions & future work



Consider multiple perturbation factors to devise more realistic metrics



•Use the metrics to measure the robustness of DLS techniques lower in the hierarchy

Simulate and use the proposed metrics to test the robustness of the DLS and ADLS techniques using an eventbased simulator

Incorporate the proposed metrics in scheduling techniques, to test their robustness when executing real world irregular scientific applications on real heterogeneous computing systems

### Summary

New complex classes of applications; uncertain computing environments; performance is more than execution time

Tradeoffs:

Numerical efficiency & stability

Scheduling overhead & load imbalance

Robustness metrics

w.r.t load variations (flexibility)

w.r.t resource failures (resilience)

Usefulness: automatic selection of robust, efficient & optimal scheduling techniques















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