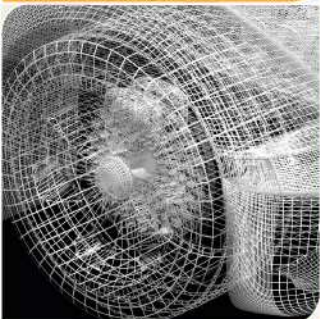


Large Scale Complex Network Analysis using the Hybrid Combination of a MapReduce Cluster and a Highly Multithreaded System



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Tech**



College of
Computing

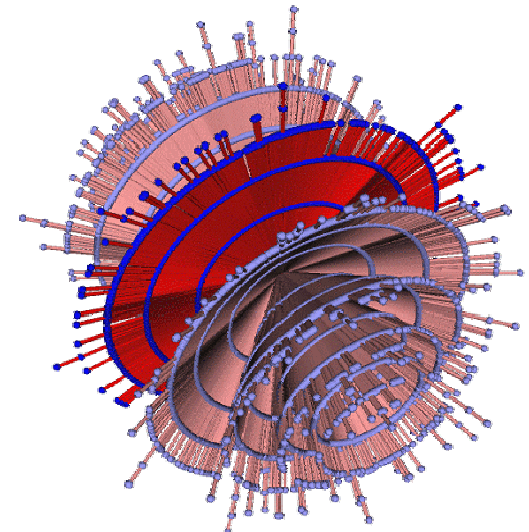
Computational Science and Engineering



A Challenge Problem

- Extracting a subgraph from a larger graph.
 - The input graph: An R-MAT* graph (undirected, unweighted) with approx. 4.29 billion vertices and 275 billion edges (7.4 TB in text format).
 - Extract subnetworks that cover 10%, 5%, and 2% of the vertices.
- Finding a single-pair shortest path (for up to 30 pairs).

| | | |
|--------|--------|-------|
| a=0.55 | a | b=0.1 |
| | c | d |
| c=0.1 | d=0.25 | |



Source: Seokhee Hong

* D. Chakrabarti, Y. Zhan, and C. Faloutsos, "R-MAT: A recursive model for graph mining," SIAM Int'l Conf. on Data Mining (SDM), 2004.



Presentation Outline

- Justify the challenge problem.
- Solve the problem using three different systems: A MapReduce cluster, a highly multithreaded system, and the hybrid system.
- Show the effectiveness of the hybrid system by
 - Algorithm level analyses
 - System level analyses
 - Experimental results



Highlights

| | A MapReduce cluster | A highly multithreaded system | A hybrid system of the two |
|-----------------------|---|---|--|
| Theory level analysis | Graph extraction: $W_{\text{MapReduce}}(n) \approx \theta(T^*(n))$ Shortest path: $W_{\text{MapReduce}}(n) > \theta(T^*(n))$ | Work optimal | Effective if $ T_{\text{hmt}} - T_{\text{MapReduce}} > n / \text{BW}_{\text{inter}}$ |
| System level analysis | Bisection bandwidth and disk I/O overhead | Limited aggregate computing power, disk capacity, and I/O bandwidth | BW_{inter} is important. |
| Experiments | Five orders of magnitude slower than the highly multithreaded system in finding a shortest path | Incapable of storing the input graph | Efficient in solving the challenge problem. |

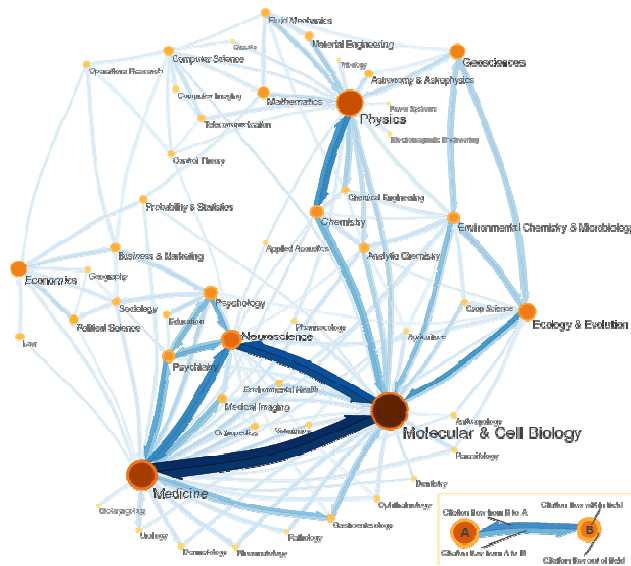


Various Complex Networks

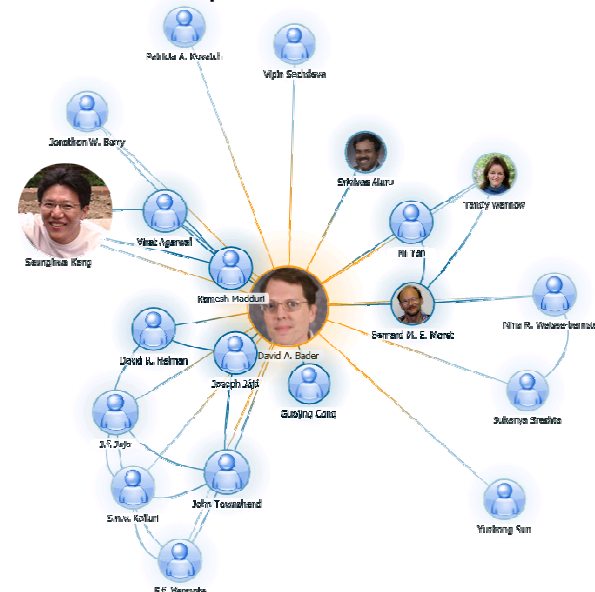
- Friendship network
- Citation network
- Web-link graph
- Collaboration network



Source: <http://www.facebook.com>



Source: <http://www.eigenfactor.org>



Source:
<http://academic.research.microsoft.com>

Extracting a graph representation from raw data



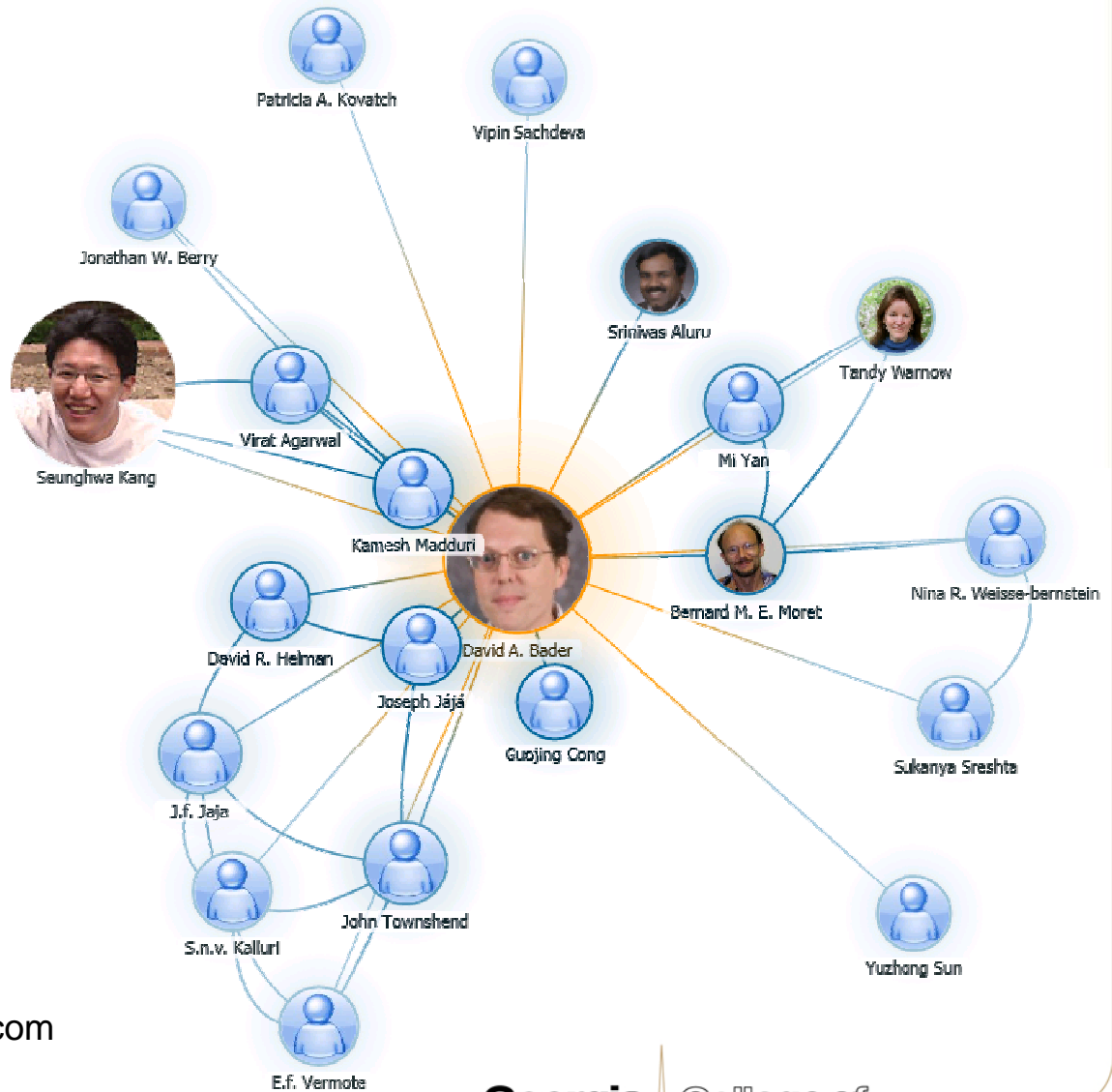
“Explore over 5,226,317 papers, 90,930 were added last week.”



Need to filter large volumes of raw data (papers) to extract a graph.

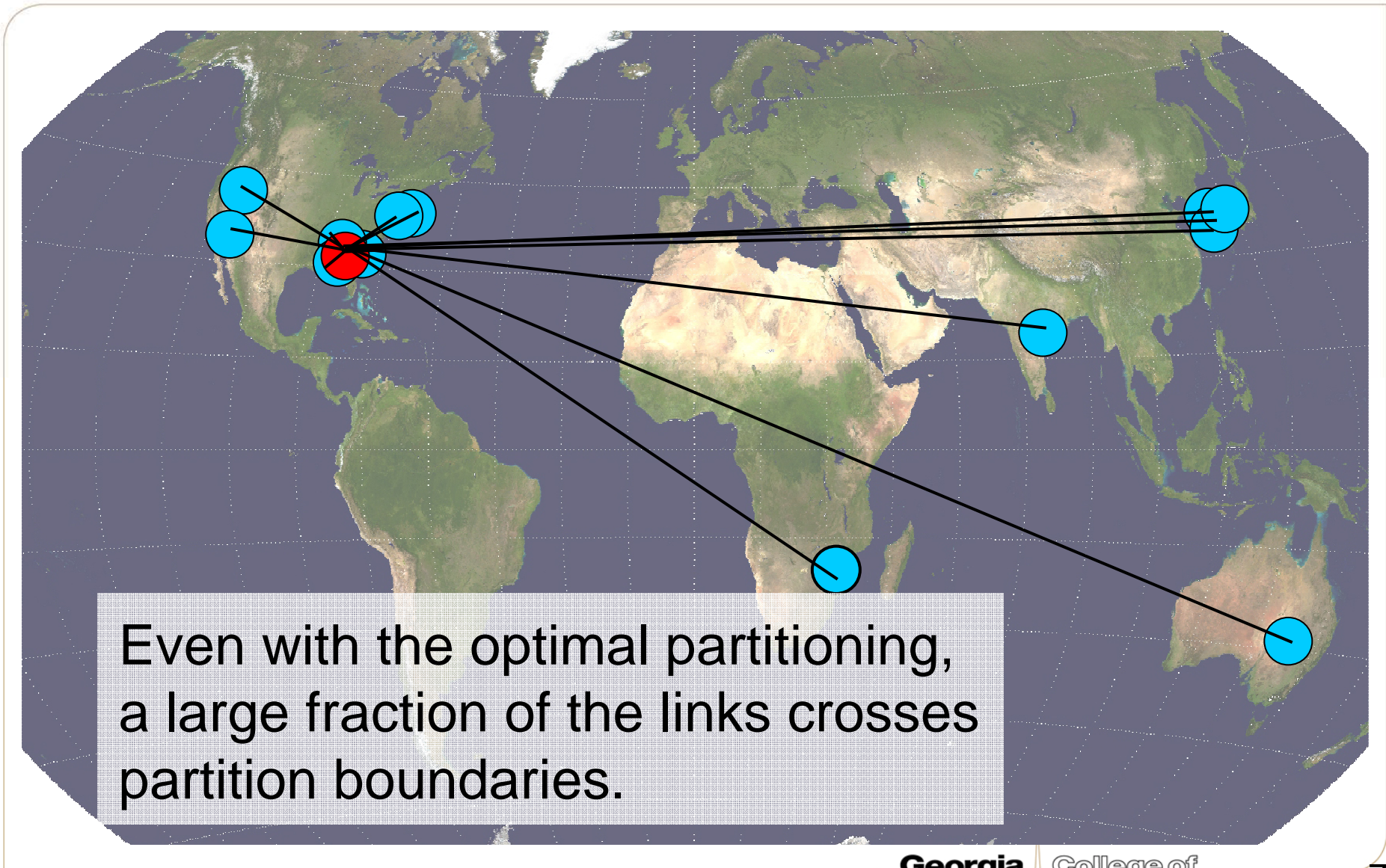
Source:

<http://academic.research.microsoft.com>





Analyzing an extracted graph

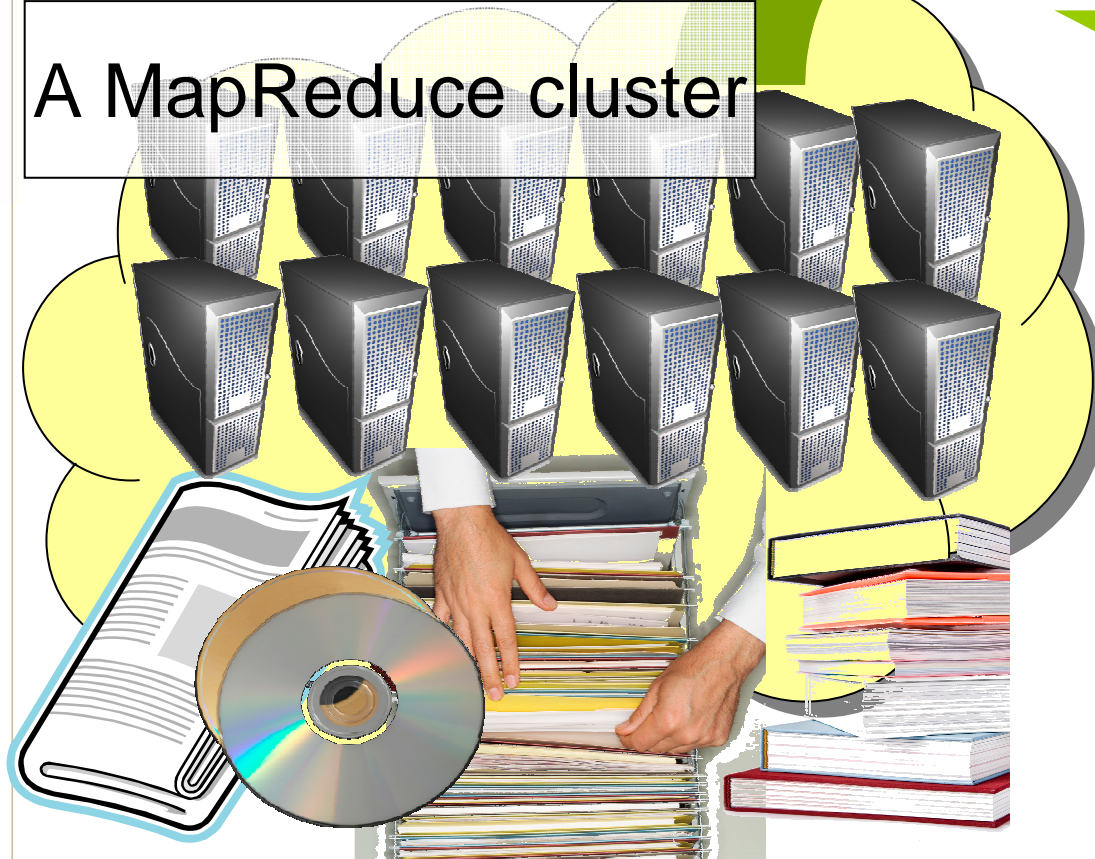


A Hybrid System to Address the Distinct Computational Challenges

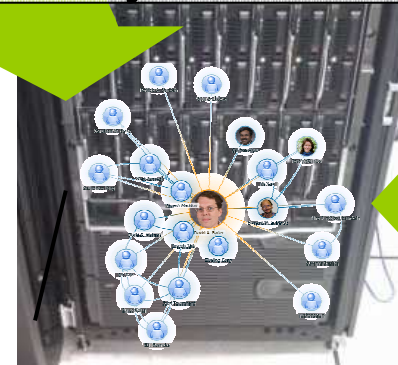


1. graph extraction

A MapReduce cluster



A highly multithreaded system

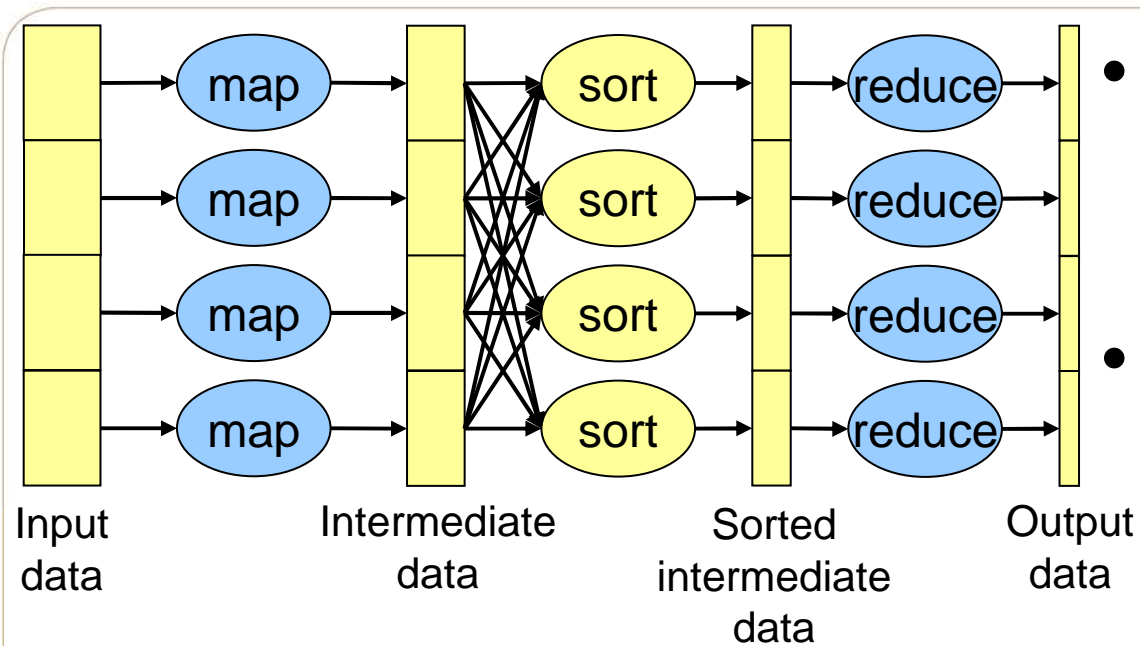


2. graph analysis queries

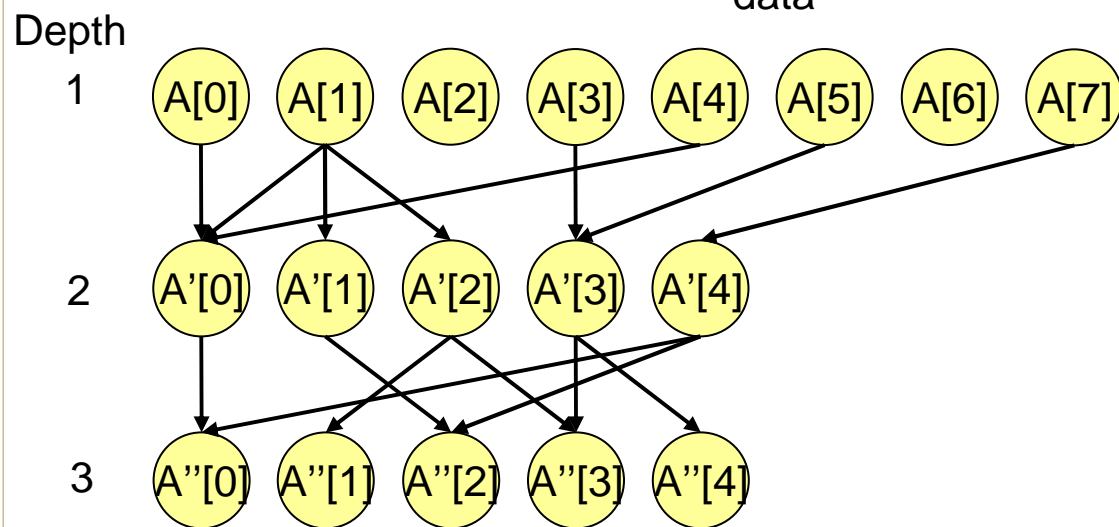




The MapReduce Programming Model



- Scans the entire input data in the map phase.
- # MapReduce iterations = the depth of a directed acyclic graph (DAG) for MapReduce computation



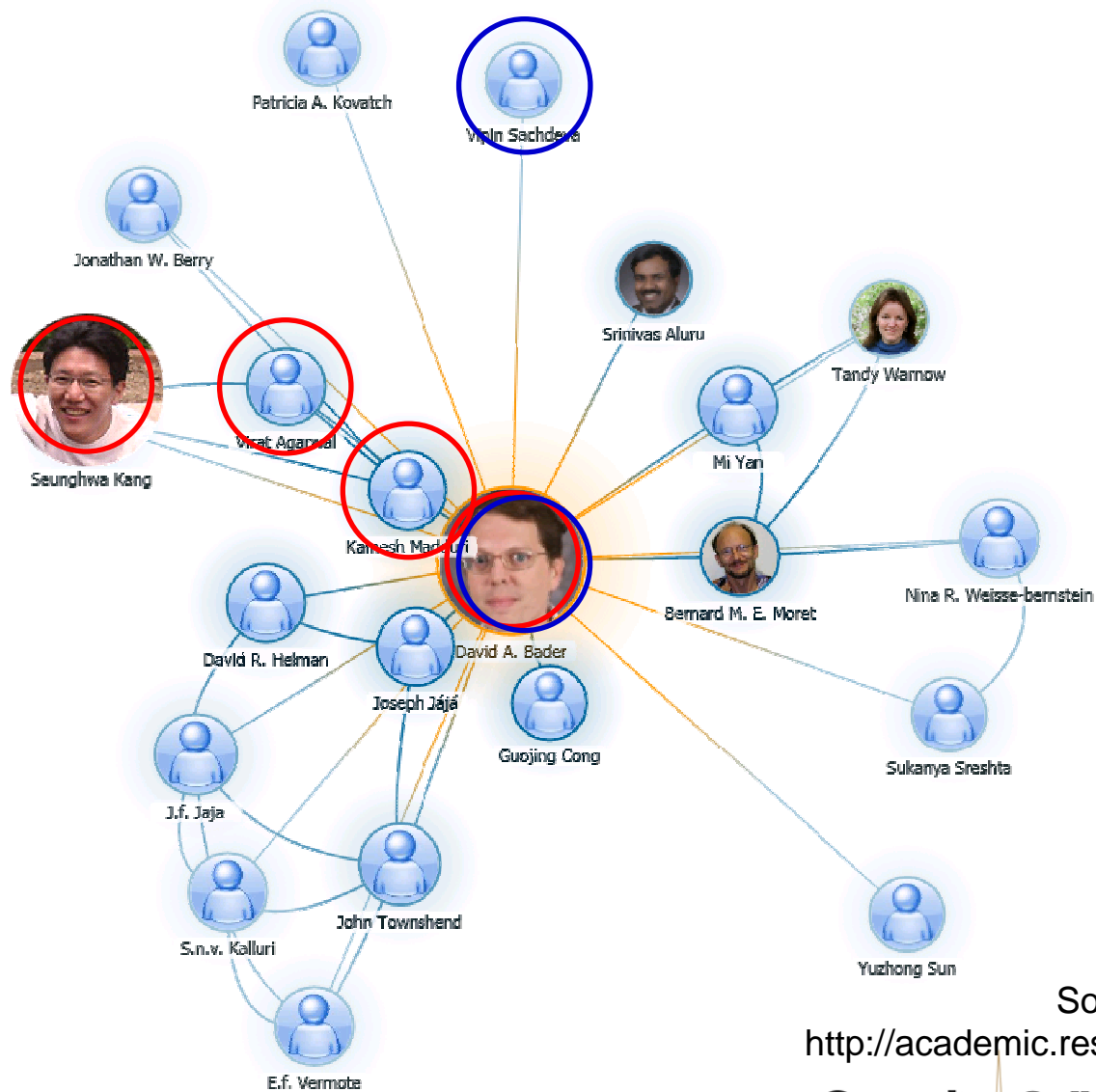
Evaluating the efficiency of MapReduce Algorithms



- $W_{\text{MapReduce}} = \sum_{i=1 \text{ to } k} (O(n_i \cdot (1 + f_i) \cdot (1 + r_i)) + p_r \cdot \text{Sort}(n_i f_i / p_r))$
 - k : # MapReduce iterations.
 - n_i : the input data size for the i th iteration.
 - f_i : map output size / map input size
 - r_i : reduce output size / reduce input size.
 - p_r : # reducers
- Extracting a subgraph
 - $k = 1$ and $f_i \ll 1 \rightarrow W_{\text{MapReduce}}(n) \approx \theta(T^*(n))$, $T^*(n)$: the time complexity of the best sequential algorithm
- Finding a single-pair shortest path
 - $k = \lceil d/2 \rceil$, $f_i \approx 1 \rightarrow W_{\text{MapReduce}}(n) > \theta(T^*(n))$



A single-pair shortest path



Source:
<http://academic.research.microsoft.com>



Bisection Bandwidth

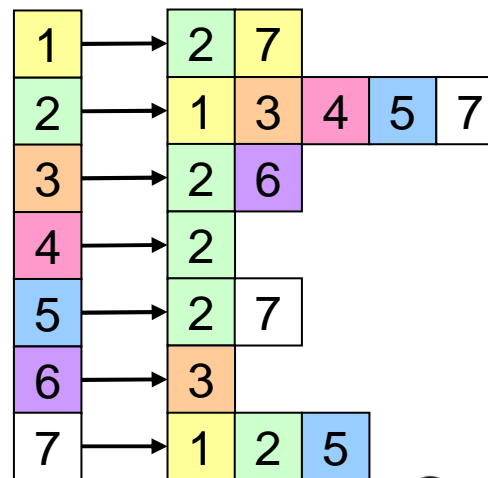
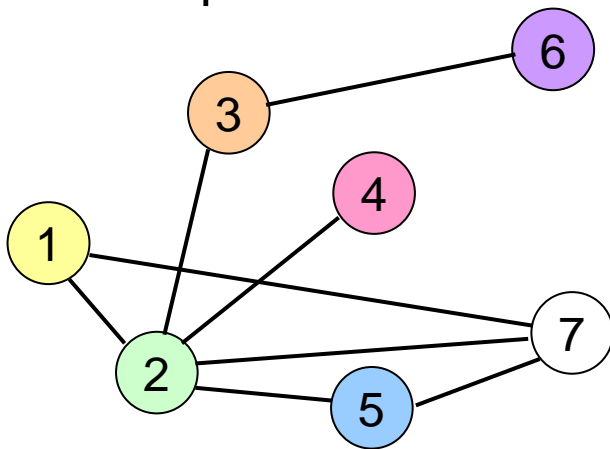
Requirements for a MapReduce Cluster

- The shuffle phase, which requires inter-node communication, can be overlapped with the map phase.
- If $T_{\text{map}} > T_{\text{shuffle}}$, T_{shuffle} does not affect the overall execution time.
 - T_{map} scales trivially.
 - To scale T_{shuffle} linearly, bisection bandwidth also needs to scale in proportion to a number of nodes. Yet, the cost to linearly scale bisection bandwidth increases super-linearly.
 - If $f \ll 1$, the sub-linear scaling of T_{shuffle} does not increase the overall execution time.
 - If $f \approx 1$, it increases the overall execution time.



Disk I/O overhead

- Disk I/O overhead is unavoidable if the size of data overflows the main memory capacity.
- Raw data can be very large.
- Extracted graphs are much smaller.
 - The Facebook network: 400 million users \times 130 friends per user \rightarrow less than 256 GB using the sparse representation.

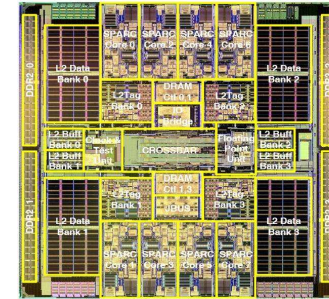


A Highly Multithreaded System w/ the Shared Memory Programming Model



- Provide a random access mechanism.
- In SMPs, non-contiguous accesses are expensive.*
- Multithreading tolerates memory access latency.+
- There is a work optimal parallel algorithm to find a single-pair shortest path.

Sun Fire T2000 (Niagara)



Features:

- Eight 64b Multithreaded SPARC Cores
- Shared 3MB L2 Cache
- 16KB ICache per Core
- 8KB DCache per Core
- Four 144b DDR-2 DRAM Interfaces (400 MT/s)
- 3.2GB/s JBUS I/O
- Crypto: Public Key (RSA)
- Extensive RAS

Technology:

- 90nm CMOS Process
- 9LM Copper Interconnect
- Power: 63 Watts @ 1.2GHz
- Die Size: 378mm²
- 279M Transistors
- Package: Flip-chip ceramic LGA (1933 pins)

Source: Sun Microsystems

Cray XMT



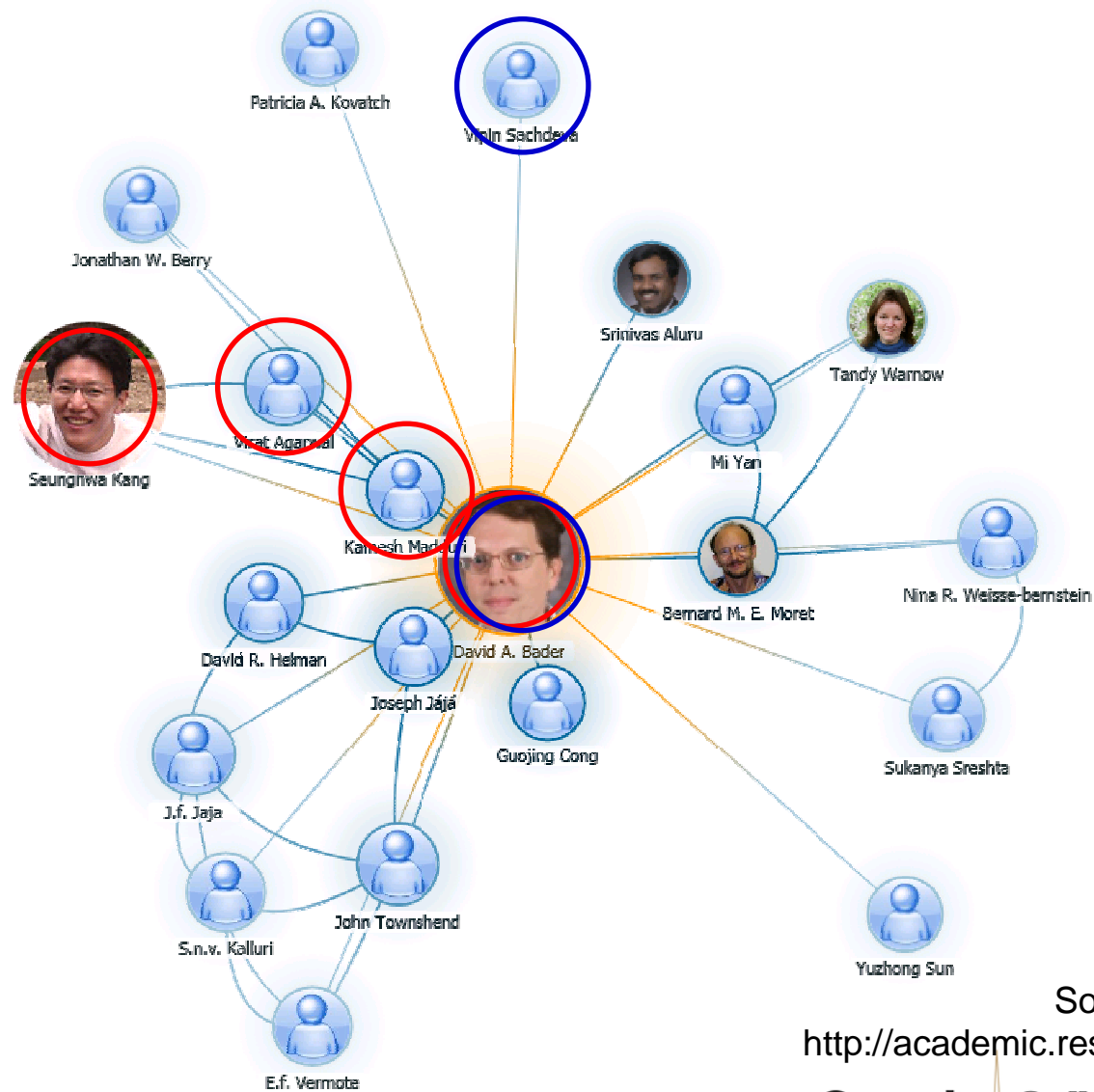
Source: Cray

* D. R. Helman and J. Ja'Ja', "Prefix computations on symmetric multiprocessors," J. of parallel and distributed computing, 61(2), 2001.

+ D. A. Bader, V. Kanade, and K. Madduri, "SWARM: A parallel programming framework for multi-core processors," Workshop on Multithreaded Architectures and Applications, 2007.



A single-pair shortest path



Source:
<http://academic.research.microsoft.com>

Low Latency High Bisection Bandwidth Interconnection Network



- Latency increases as the size of a system increases.
 - A larger number of threads and additional parallelism are required as latency increases.
- Network cost to linearly scale bisection bandwidth increases super-linearly.
 - But not too expensive for a small number of nodes.
- These limit the size of a system.
 - Reveal limitations in extracting a subgraph from a very large graph.



The Time Complexity of an Algorithm on the Hybrid System

- $T_{\text{hybrid}} = \sum_{i=1 \text{ to } k} \min(T_{i, \text{MapReduce}} + \Delta, T_{i, \text{hmt}} + \Delta)$
 - k : # steps
 - $T_{i, \text{MapReduce}}$ and $T_{i, \text{hmt}}$: time complexities of the i_{th} step on a MapReduce cluster and a highly multithreaded system, respectively.
 - Δ : $n_i / \text{BW}_{\text{inter}} \times \delta(i-1, i)$,
 - n_i : the input data size for the i_{th} step.
 - BW_{inter} : the bandwidth between a MapReduce cluster and a highly multithreaded system.
 - $\delta(i-1, i)$: 0 if selected platforms for the $i-1_{\text{th}}$ and i_{th} steps are same. 1, otherwise.



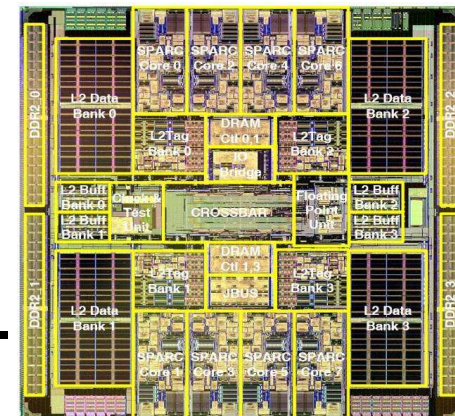
Test Platforms

- A MapReduce cluster
 - 4 nodes
 - 4 dual core 2.4 GHz Opteron processors and 8 GB main memory per node.
 - 96 disks (1 TB per disk).
- A highly multithreaded system
 - A single socket UltraSparc T2 1.2 GHz processor (8 core, 64 threads).
 - 32 GB main memory.
 - 2 disks (145 GB per disk)
- A hybrid system of the two



Source: <http://hadoop.apache.org/>

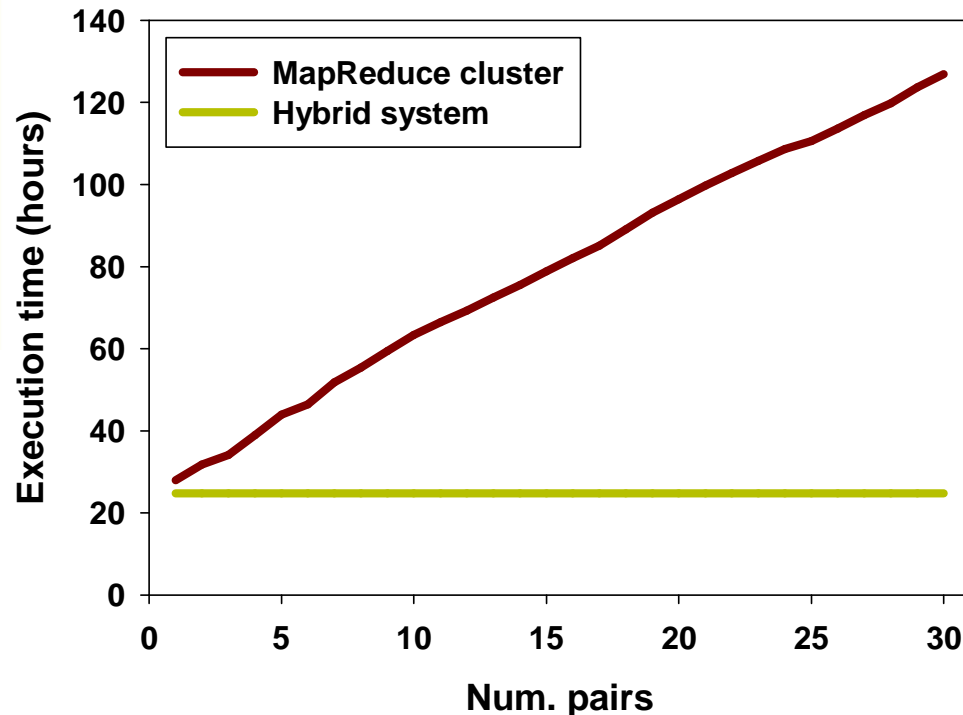
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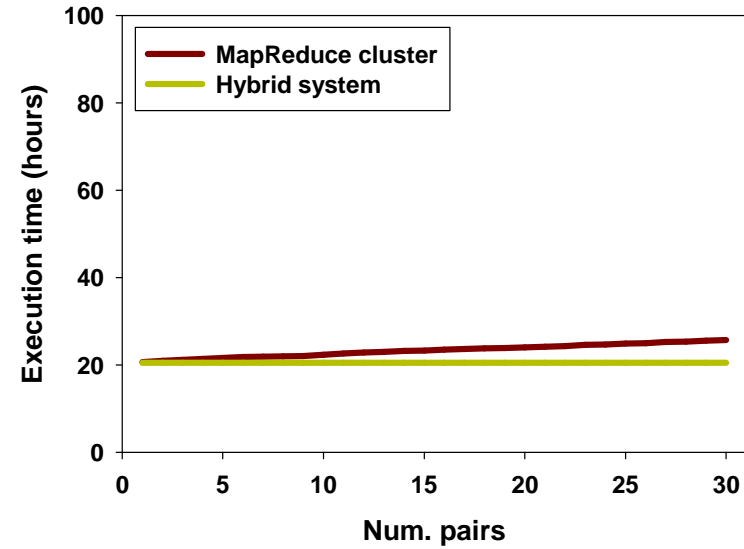
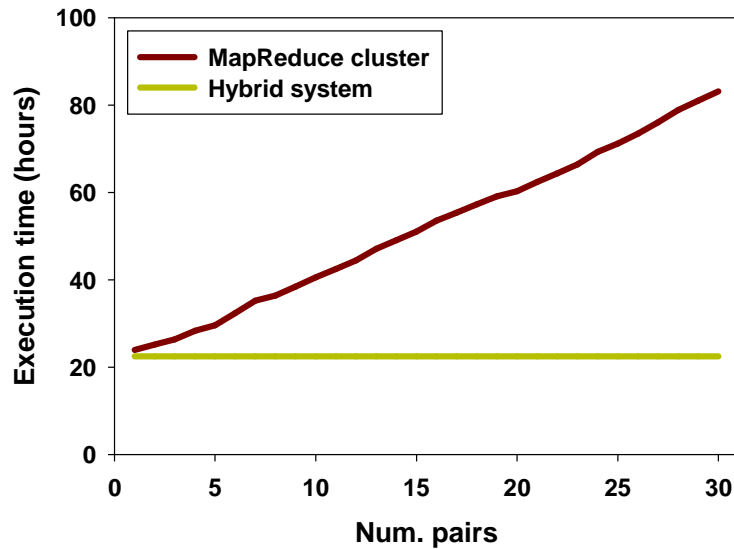
A subgraph that covers 10% of the input graph



| | MapReduce | Hybrid |
|--|-----------|---------|
| Subgraph extraction | 24 | 24 |
| Memory loading | - | 0.83 |
| Finding a shortest path (for 30 pairs) | 103 | 0.00073 |

Once the subgraph is loaded into the memory, the hybrid system analyzes the subgraph five orders of magnitude faster than the MapReduce cluster (103 hours vs 2.6 seconds).

Subgraphs that cover 5% (left) and 2% (right) of the input graph



| | MapReduce | Hybrid |
|--|-----------|---------|
| Subgraph extraction | 22 | 22 |
| Memory loading | - | 0.42 |
| Finding a shortest path (for 30 pairs) | 61 | 0.00047 |

| | MapReduce | Hybrid |
|--|-----------|---------|
| Subgraph extraction | 21 | 21 |
| Memory loading | - | 0.038 |
| Finding a shortest path (for 30 pairs) | 5.2 | 0.00019 |



Conclusions

- Performance and programmability are highly correlated with the match between a workload's computational requirements and a programming model and an architecture.
- Our hybrid system is effective in addressing the distinct computational challenges in large scale complex network analysis.



Acknowledgment of Support

