



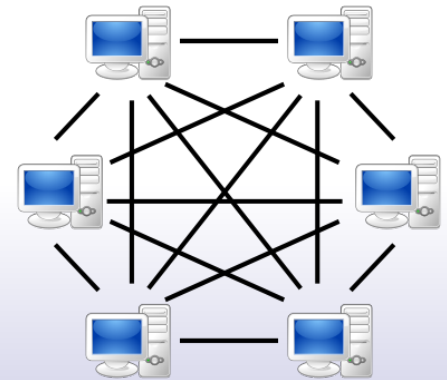
BitTorrent and fountain codes: friends or foes?

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Peer to peer and file sharing applications

Peer to peer paradigm has a huge diffusion nowadays [1]



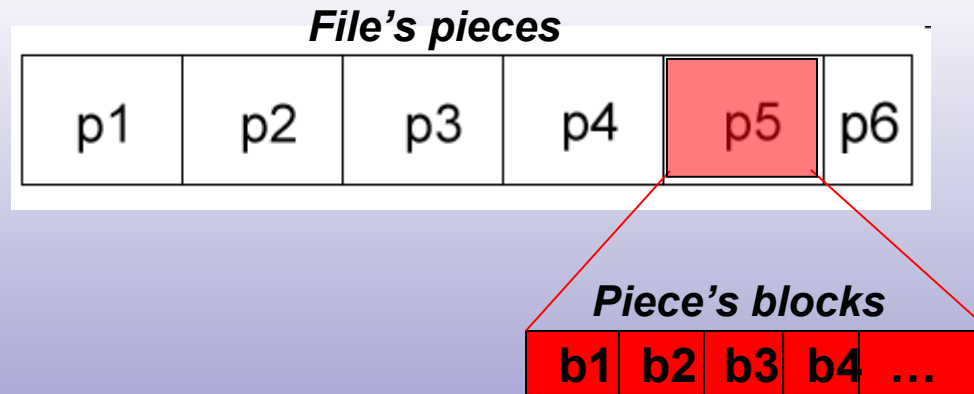
It can be employed for a lot of different applications:

- **File sharing:** BitTorrent [2][3] , eDonkey, eMule, DC++
- **Video Streaming :** SopCast, PPLive, PPStream
- **Distributed portals:** Osiris
- **Others:** Skype, Sciencenet, Spotify

BitTorrent is doubtless the reference architecture for file sharing



BitTorrent adopts a **multi-part** download **scheme**:



- The file is divided into **pieces**, each piece into **blocks**.
- Peers are **synchronized** between them **at piece level**.
- Each client knows what pieces have been **completely downloaded** from his neighbors.
- Peers exchange between them piece's blocks



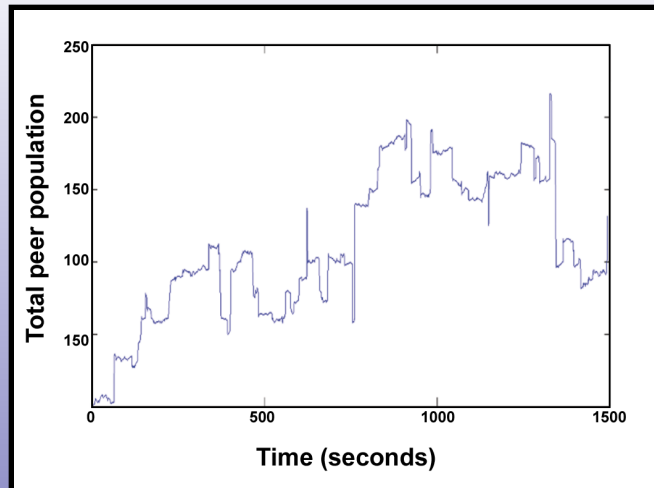
Main **BitTorrent's** strategies:

- **Tit-for-Tat**: assures reciprocity between downloading and uploading rates.
- **Rarest first**: assures a fair distribution of file's pieces
Note that only complete pieces can be shared and distributed

Many research works claim that BitTorrent has **near-optimal** performance [4] [5] [6]



Some phenomena can **cripple** BitTorrent performance:



Example of simulated population's dynamics in a network affected by flash crowd and high Churn rates.

- **Flash crowd:** many peer join or leave the network at the same time
- **High churn rates:** peers join and leave the network at high rates
- Many peer **leave** the overlay network **at the same time**



Modifying BitTorrent protocol 1: Luby-transform codes

LT codes [7][8] are a class a *fountain codes* that use simple XOR operations to encode and decode the message. The decoding process has a certain overhead ϵ : the recipient needs to receive a number of packets that exceed in percentage the original size of the message by a certain amount.

If k is the original number of non coded blocks, the decoding process succeeds when $k(1+ \epsilon)$ packets are received.

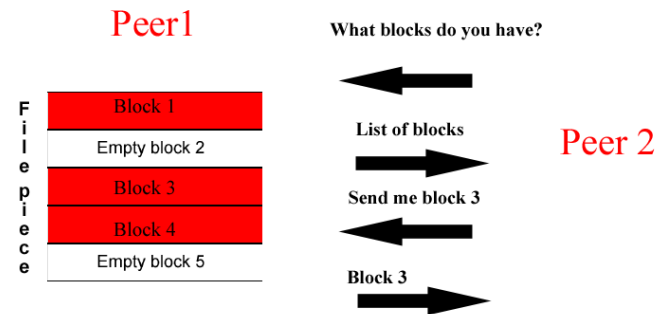
Benefits

- Avoid **content reconciliation**
- A peer can share **partially downloaded** file parts
- File parts can be download **concurrently** from different peers

Drawbacks

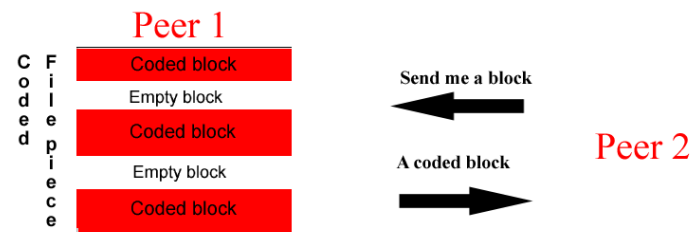
- Decoding **overhead**
- Encoder/decoder complexity

Uncoded protocol



But block 3 is only 16KB: is not convenient share partial blocks, **too much signalling overhead**

Using LT codes



Coded blocks can be sent without content reconciliation



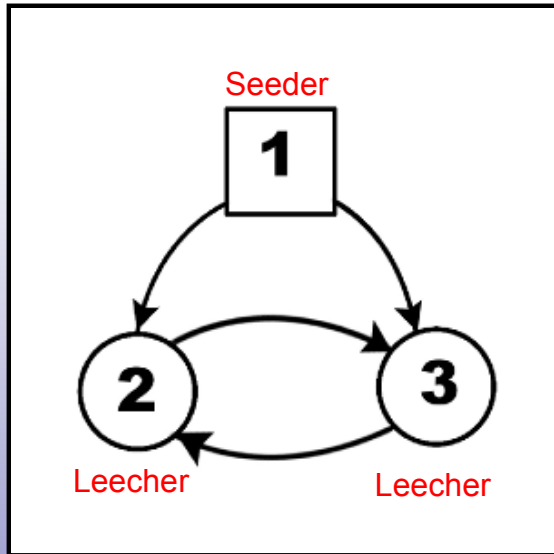
Summary of protocol modifications:

- Use **LT codes** to code file's pieces. The file's piece subdivision is different in accordance with the requirements of encoding. The decoding overhead is about 10%
- Share **partial** file's piece
- A piece can be downloaded from **more peers at same time** without content reconciliation
- The protocol requests first **the rarest complete piece**. If there are not complete pieces available, a peer requests the **rarest partial piece**
- If a piece is fully decoded a peer sends **new coded blocks**

We use **General purpose simulator GPS**^[9] simulator in order to test our modified protocol



Experimental results: simple three peer topology



A simple topology with a seeder and two leechers

DOWNLOADING TIMES AND ACHIEVABLE BITRATE OBSERVED ON A SIMPLE TOPOLOGY OF THREE PEERS.

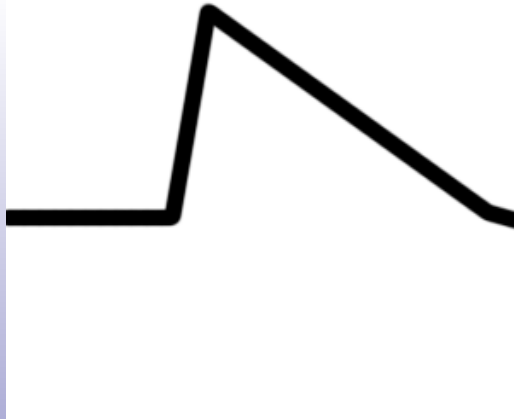
File size	BTd ¹	LTd ²	BTb ¹	LTb ²	Gain (%)
1 MBytes	38s	35s	210Kbps	228Kbps	8%
2 MBytes	50s	48s	320Kbps	333Kbps	4%
4 MBytes	77s	75s	415Kbps	426Kbps	2.6%

¹ the columns **BTd** and **BTb** show respectively the download times and achievable bitrate of the standart BitTorrent protocol

² the columns **LTd** and **LTb** show respectively the download times and achievable bitrate of the modified protocol that use the LT codes



Single flash crowd scenario set up:



- Time **0**: a flash crowd of 50 peers occurs
- Time **50**: 20 random selected peers leave the network, regardless of the state of the download

DOWNLOADING TIMES AND ACHIEVABLE BITRATE OBSERVED IN A SINGLE FLASH CROWD SCENARIO.

File size	PS ¹	BTd ²	LTd ³	BTb ²	LTb ³	Gain
1 MByte	S	53.1s	18.1s	151 Kbps	441 Kbps	66%
1 MByte	M	23.9s	17.1s	334 Kbps	467 Kbps	26%

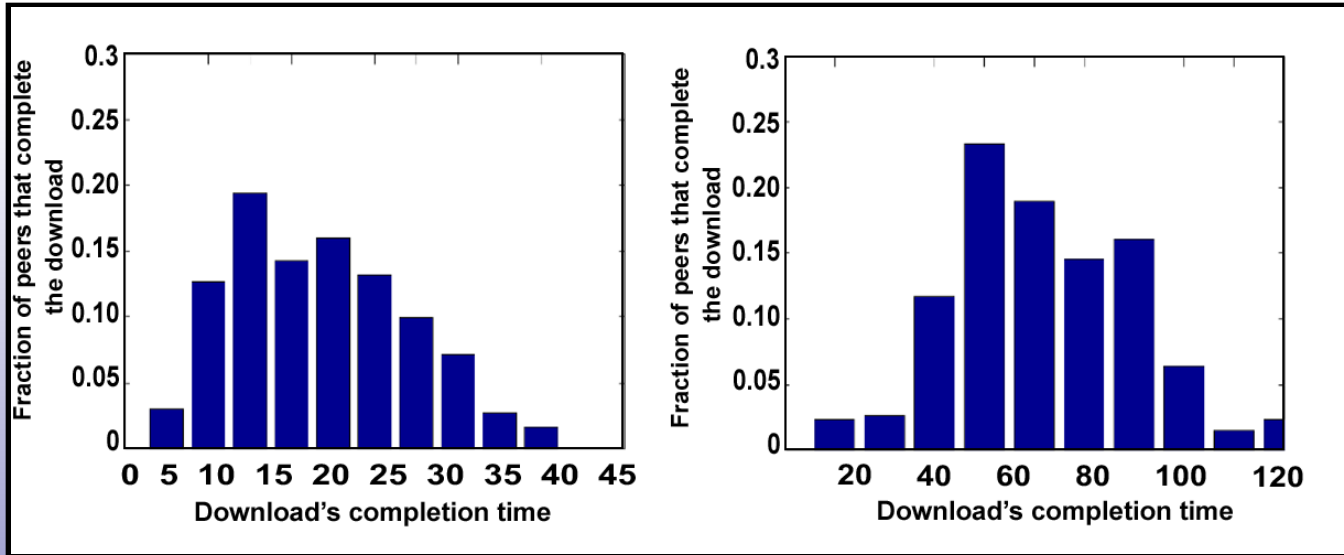
¹ the file subdivision, single (S) or multiple (M) pieces

² the columns **BTd** and **BTb** show respectively the download times and achievable bitrate of the standart BitTorrent protocol

³ the columns **LTd** and **LTb** show respectively the download times and achievable bitrate of the modified protocol that use the LT codes



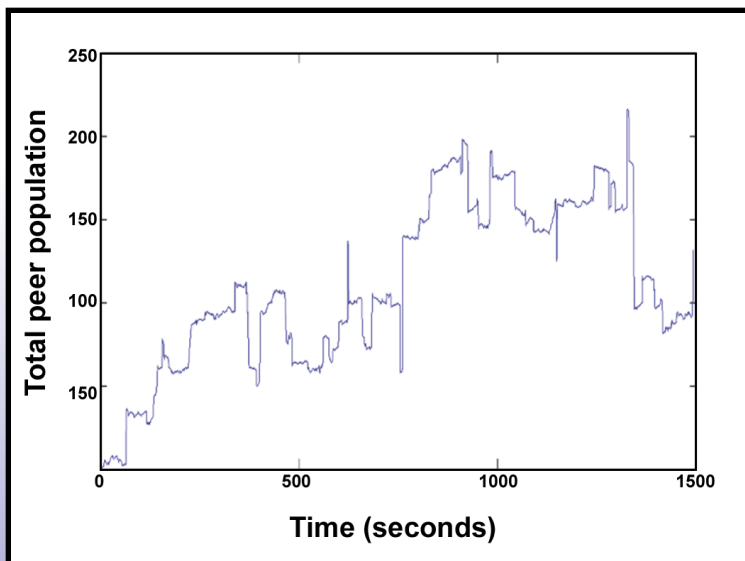
Experimental results: downloading times distribution



Comparison of downloading time distributions between original protocol and the modified one



Experimental results: a more complex scenario 1



During the simulation multiple flash crowd and massive departure of peer occur

Example of simulated population's dynamics in a network affected by flash crowd and high Churn rates.

DOWNLOADING TIMES AND ACHIEVABLE BITRATE OBSERVED IN A SCENARIO AFFECTED BY MULTIPLE FLASH CROWD AND MASSIVE DEPARTURE OF PEERS.

File size	BTd ¹	LTd ²	BTb ¹	LTb ²	Gain
4 MBytes	65 s	56 s	492 Kbps	571 Kbps	14%
8 MBytes	145 s	129 s	441 Kbps	496 Kbps	12%

¹ the columns **BTd** and **BTb** show respectively the download times and achievable bitrate of the standard BitTorrent protocol

² the columns **LTd** and **LTb** show respectively the download times and achievable bitrate of the modified protocol that use the LT codes



In this paper we proposed a novel modifications of the **BitTorrent protocol** by introducing the **LT encoding**. Then we proved by simulations that better performance can be achieved considering **file of small size** and in **adverse network conditions**. In such situation our proposed protocol yields a gain that is above the 10% in all simulated scenarios, in spite of the decoding overhead that the LT codes introduce.



Thank you for your kind attention



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- [2] <http://www.bittorrent.com/>
- [3] BitTorrent protocol specification v1.0 <http://wiki.theory.org/BitTorrentSpecification>, June 2005
- [4] A. Legout, G. Urvoy-Keller, P. Michiardi, *Understanding bittorrent: An experimental perspective. Technical report*, CNRS : FRE2660 EURECOM
- [5] A. Legout, G. Urvoy-Keller, P. Michiardi, *Rarest First and Choke Algorithms Are Enough*, IMC '06: Proceedings of the 6th ACM SIGCOMM conference on Internet measurement, Rio de Janeiro, Brazil, 2006.
- [6] A. Al-Hamra, A. Legout, and C. Barakat, *Understanding the Properties of the BitTorrent Overlay*, CoRR abs/0707.1820: (2007).
- [7] M. Luby, *LT codes*, Foundations of Computer Science, 2002. Proceedings. The 43rd Annual IEEE Symposium on (2002).
- [8] E. Hyttia, T. Tirronen, J. Virtamo *Optimizing the Degree Distribution of LT Codes*, in RESIM 2006, 6th International Workshop on Rare Event Simulation, 2006, Bamberg, Germany
- [9] W. Yang, N. Abu-Ghazaleh, *GPS: A General Peer-to-Peer Simulator and its Use for Modeling BitTorrent*, Proceedings of 13th Annual Meeting of the IEEE International Symposium on Modeling, Analysis, and Simulation of Computer and Telecommunication Systems (MASCOTS '05) , Sept 27-29, 2005, Atlanta.